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A NOTE ON INTUITIONISTIC SENTENTIAL CALCULUS (*ISC*)

Let FM be the set of all formulas of a sentential language with usual connectives $\neg, \wedge, \vee, \Rightarrow$ and \Leftrightarrow and, let IA be the set of logical axioms of *ISC*, see [1]. Every subset of FM containing IA and closed under the modus ponens rule is called a theory. *TAUT* is the smallest theory. The quotient of FM modulo *TAUT* is the free pseudo-Boolean algebra. By the prime filter theorem for distributive lattices: if α is not in the theory T then there exists a prime theory T_0 over T such that α is not in T_0 .

Let TV be the set of all t in 2^{FM} such that for all α, β in FM :

- (0) $t(\alpha) = 1$ whenever α is in IA
- (1) $t(\alpha) = 0$ or $t(\neg\alpha) = 0$
- (2) $t(\alpha \wedge \beta) = 1$ iff $t(\alpha) = t(\beta) = 1$
- (3) $t(\alpha \vee \beta) = 0$ iff $t(\alpha) = t(\beta) = 0$
- (4) either $t(\alpha \Rightarrow \beta) = 0$ or $t(\alpha) = 0$ or $t(\beta) = 1$
- (5) $t(\alpha \Leftrightarrow \beta) = 1$ iff $t(\alpha \Rightarrow \beta) = t(\beta \Rightarrow \alpha) = 1$

Obviously, if T is any theory then $t(\alpha) = 1$ for every α in T and each t in TV . On the other hand, if α is not in the theory T then there exists t in TV such that $t(T) = 1$ and $t(\alpha) = 0$.

The intuitionistic consequence relation is defined for any subset X of FM and each α in FM as follows:

$X \vdash \alpha$ iff α belongs to every theory over X .

THEOREM. $X \vdash \alpha$ iff for all t in TV , $t(\alpha) = 1$ whenever $t(X) = 1$.

COROLLARY. The set *TAUT* is decidable.

COMMENT. Dr Basil Discord once told me that the *ISC* is a (logically) two-valued logic with truth-functional conjunction and disjunction and, non-truth-functional negation, implication and equivalence.

References

- [1] H. Rasiowa and R. Sikorski, **The Mathematics of the Meta-mathematics**, PWN, Warszawa 1963.

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